

Lonely Points in Simplices

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Abstract Given a lattice $L \subseteq \mathbb{Z}^m$ and a subset $A \subseteq \mathbb{R}^m$, we say that a point in A is *lonely* if it is not equivalent modulo L to another point of A . We are interested in identifying lonely points for specific choices of L when A is a dilated standard simplex, and in conditions on L which ensure that the number of lonely points is unbounded as the simplex dilation goes to infinity.

Keywords integer points, polytopes, lattices, discrete geometry

1 Introduction

The geometric problem considered in this article arose from an attempt to construct an algorithm for simplifying so-called C-finite sequences. A sequence $(a_n)_{n=0}^{\infty}$ in the field \mathbb{C} of complex numbers is called *C-finite* [11] if it satisfies a linear recurrence with constant coefficients¹, i.e., if there are constants

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¹ W.l.o.g. we consider \mathbb{C} instead of an algebraically closed arbitrary field of characteristic zero

$c_0, \dots, c_r \in \mathbb{C}$, not all zero, such that

$$c_0 a_n + c_1 a_{n+1} + \dots + c_r a_{n+r} = 0$$

for all $n \in \mathbb{N}$. A standard example is the sequence of Fibonacci numbers (take $c_0 = c_1 = 1$ and $c_2 = -1$). C-finite sequences and their properties are very well understood [16, 7, 11, 17, 12]. In particular, it is known that a sequence is C-finite if and only if it can be expressed as a linear combination of exponential terms with polynomial coefficients, i.e., if there are polynomials $p_1, \dots, p_m \in \mathbb{C}[x]$ and constants $\phi_1, \dots, \phi_m \in \mathbb{C}$, such that

$$a_n = p_1(n)\phi_1^n + \dots + p_m(n)\phi_m^n$$

for all $n \in \mathbb{N}$. If the ϕ_i are pairwise distinct and all the p_i are nonzero, then the order r of the corresponding recurrence turns out to be $m + \sum_{i=1}^m \deg(p_i)$.

One of the consequences of the characterization above is that the class of C-finite sequences is closed under addition and multiplication, i.e., when the sequences $(a_n)_{n=0}^\infty$ and $(b_n)_{n=0}^\infty$ are C-finite, then so are the sequences $(a_n + b_n)_{n=0}^\infty$ and $(a_n b_n)_{n=0}^\infty$. In particular, when we plug a C-finite sequence into a polynomial, the result is again a C-finite sequence. For example, since the sequence $(F_n)_{n=0}^\infty$ of Fibonacci-numbers is C-finite, so is the sequence $(5F_n^3 - 7F_n^2 + 9F_n - 4)_{n=0}^\infty$ obtained by plugging $(F_n)_{n=0}^\infty$ into the polynomial $5x^3 - 7x^2 + 9x - 4 \in \mathbb{C}[x]$.

Given a C-finite sequence $(a_n)_{n=0}^\infty$, specified by a recurrence of order r and a set of initial values, we want to decide whether there is a polynomial $q \in \mathbb{C}[x]$ of positive degree such that the C-finite sequence $(q(a_n))_{n=0}^\infty$ satisfies a recurrence of order less than r . This problem is of interest because certain number-theoretic questions about C-finite sequences can at the moment only be answered when the recurrence order is not too large, e.g. see [14]. By using results of our paper to pass from $(a_n)_{n=0}^\infty$ to $(q(a_n))_{n=0}^\infty$, we hope to extend the scope of these algorithms and advance, for example, their use in applications of static analysis of computer systems, see for example [10, 9].

The construction of an algorithm for finding $q \in \mathbb{C}[x]$, such that $(q(a_n))_{n=0}^\infty$ yields a C-finite sequence of lower order than a , has led us to the following geometric problem. Let $S \subseteq \mathbb{R}^m$ be the standard simplex, i.e., the convex hull of 0 and the unit vectors $e_1, \dots, e_m \in \mathbb{R}^m$. Moreover, let $L \subseteq \mathbb{Z}^m$ be a lattice, i.e., an additive subgroup of \mathbb{Z}^m . Two points $u, v \in \mathbb{R}^m$ are called equivalent modulo L if we have $u - v \in L$. We consider the integer points in a dilation dS of S , for some $d > 0$. A point $u \in dS \cap \mathbb{Z}^m$ is called *lonely* if there does not exist any other point $v \in dS \cap \mathbb{Z}^m$ such that $u - v \in L$. In this paper, we are interested to describe properties of these lonely points.

In Section 2, we will give some more details on how the original problem about C-finite sequences leads to the consideration of lonely points. This material is provided only as background information and not strictly needed for the rest of the paper. In Section 3, we summarize basic definitions and facts about cones, simplices, and lattices, and fix the notation we use. In Section 4 we present algorithms that for a given lattice L and a given d determine all

the lonely points, and recognize whether the number is unbounded as d goes to infinity. Finally, in Section 5 we derive a sufficient condition on the lattice that guarantee that the number of lonely points is unbounded.

2 Ansatz and Exponent Lattice

Consider a \mathbb{C} -finite sequence $(a_n)_{n=0}^{\infty}$ which satisfies a recurrence of order r . We want to know whether there is a polynomial $q \in \mathbb{C}[x] \setminus \mathbb{C}$ such that $(q(a_n))_{n=0}^{\infty}$ satisfies a recurrence of lower order. If we have an upper bound d on the degree of q , then this question can be answered as follows:

1. Compute $p_1, \dots, p_m \in \mathbb{C}[x]$ and $\phi_1, \dots, \phi_m \in \mathbb{C}$ such that $a_n = p_1(n)\phi_1^n + \dots + p_m(n)\phi_m^n$ for all $n \in \mathbb{N}$ (see [11] for how to do this).
2. Make an ansatz $q = q_0 + q_1x + \dots + q_dx^d$ with undetermined coefficients q_0, \dots, q_d , plug the closed form representation of step 1 into p .
3. Write the resulting expression in the form

$$\bigcirc \psi_1^n + \dots + \bigcirc \psi_\ell^n$$

where the $\psi_i \in \mathbb{C}$ are pairwise distinct and the \bigcirc are polynomials in n whose coefficients are \mathbb{C} -linear combinations of the unknowns q_0, \dots, q_d .

4. For every subset $I \subseteq \{1, \dots, \ell\}$ such that $|\{1, \dots, \ell\} \setminus I| = r - 1$, equate the coefficients with respect to n in all the \bigcirc expressions belonging to some ψ_i with $i \in I$ to zero and solve the resulting linear system for the unknowns q_0, \dots, q_d . If the solution space contains a vector (q_0, \dots, q_d) in which not only q_0 is nonzero, return the corresponding polynomial $q_0 + q_1x + \dots + q_dx^d$. Otherwise, try the next I .
5. When no subset I yields a solution, return “there is no such q ”.

- Example 1*
1. The \mathbb{C} -finite sequence $(a_n)_{n=0}^{\infty}$ with $a_n = 1 + 2^n + 2^{-n}$ satisfies a recurrence of order 3 and no lower order recurrence. With $d = 2$, the algorithm sketched above finds the polynomial $q(x) = x^2 - 2x - 1$. Indeed, $q(a_n) = 4^n + 4^{-n}$ satisfies a recurrence of order 2.
 2. The \mathbb{C} -finite sequence $(a_n)_{n=0}^{\infty}$ with $a_n = 1 + 3^n + 3^{2n} + 2 \cdot 3^{3n} - 2 \cdot 3^{4n}$ satisfies a recurrence of order 5 and no lower order recurrence. For this input, the algorithm finds the polynomial $q(x) = x^2 - 3x + 2$, and indeed, $q(a_n) = -3^n + 7 \cdot 3^{4n} - 8 \cdot 3^{7n} + 4 \cdot 3^{8n}$ satisfies a recurrence of order 4. Similar examples can be constructed using polynomials with sparse powers. Such polynomials have been studied for example in [5].
 3. The \mathbb{C} -finite sequence $(a_n)_{n=0}^{\infty}$ with $a_n = 1 + 2^n - 2^{-n}$ satisfies a recurrence of order 3, and with the algorithm sketched above we can show that there is no polynomial q of degree $d \leq 5$ such that $q(a_n)$ satisfies a recurrence of order 2.

When we have checked the existence of a polynomial q for a specific degree d and found that no such polynomial exists, we can try again with a larger choice of d . It would be good to know when we can stop: starting from the

recurrence of $(a_n)_{n=0}^\infty$, can we determine a finite bound on the degree of the polynomials q that may lead to lower order recurrences?

In order to see from where such a bound could emerge, restrict the search to polynomials q with $q_d = 1$. Observe what happens in step 2 of the procedure sketched above. Plugging the expression $p_1(n)\phi_1^n + \dots + p_m(n)\phi_m^n$ into the ansatz for q produces

$$\begin{aligned} q_0 + q_1 \sum_{i=1}^m p_i(n)\phi_i^n + q_2 \sum_{i,j=1}^n p_i(n)p_j(n)(\phi_i\phi_j)^n \\ + \dots + \sum_{i_1, \dots, i_d=1}^n \prod_{j=1}^d p_{i_j}(n) \left(\prod_{j=1}^d \phi_{i_j} \right)^n, \end{aligned} \quad (1)$$

so the ψ_i 's appearing in step 3 are precisely the products $\phi_1^{v_1} \dots \phi_m^{v_m}$ with $v_1 + \dots + v_m \leq d$. If these products are all distinct, then there is no way for the above expression to vanish identically. More generally, a necessary condition for the above expression to vanish identically for some choice of q_0, \dots, q_{d-1} , not all zero, is that a sufficient amount of cancellation takes place among the various exponential sequences $((\phi_1^{v_1} \dots \phi_m^{v_m})^n)_{n=0}^\infty$.

This leads to the consideration of the so-called *exponent lattice*

$$L = \{ (v_1, \dots, v_m) \in \mathbb{Z}^m : \phi_1^{v_1} \dots \phi_m^{v_m} = 1 \} \subseteq \mathbb{Z}^m,$$

which also plays an important role for determining the algebraic relations among C-finite sequences [13]. For example, for the Fibonacci numbers, where we have $\phi_1 = \frac{1}{2}(1 + \sqrt{5})$ and $\phi_2 = \frac{1}{2}(1 - \sqrt{5})$, the exponent lattice is generated by $(2, 2)$.

A term $(\phi_1^{v_1} \dots \phi_m^{v_m})^n$ appearing in (1) cannot be canceled unless there is some other point $(\tilde{v}_1, \dots, \tilde{v}_m) \in \mathbb{N}^m$ with $\tilde{v}_1 + \dots + \tilde{v}_m \leq d$ and $(v_1 - \tilde{v}_1, \dots, v_m - \tilde{v}_m) \in L$. If d is such that r or more of the terms have no partner for cancellation, then it is clear that there is no solution q of degree d . Moreover, if L is such that the number of terms without partner tends to infinity as d increases, then there is a finite bound on the degree that a solution q may have.

3 Lattices and Cones

We start by recalling some basic concepts from discrete geometry. Further background can be found in [2], for example.

Definition 1 (Lattices) A set $L \subset \mathbb{Z}^m$ is called a *lattice* if it contains the origin and for all $u, v \in L$ and all $\alpha, \beta \in \mathbb{Z}$ also $\alpha u + \beta v$ is an element of L . For vectors $\ell_1, \dots, \ell_k \in \mathbb{Z}^m$ we write $\langle \ell_1, \dots, \ell_k \rangle$ for the smallest lattice containing ℓ_1, \dots, ℓ_k , which we call generators of the lattice. The dimension $\dim(L)$ of a lattice is defined as the dimension of the \mathbb{R} -vector space it generates.

We always view a lattice $L \subseteq \mathbb{Z}^m$ as a set of points in the ambient space \mathbb{R}^m , spanned by the unit vectors e_1, \dots, e_m . In addition, it will be convenient to let e_0 be the zero vector.

Example 2 The vectors $(3, 3)$ and $(6, 1)$ span a lattice in \mathbb{R}^2 of dimension 2. Some points in the lattice in the positive quadrant are depicted in Figure 1a. The 2-dimensional lattice spanned by the vectors $(2, 1, 0)$ and $(0, 2, 1)$ in \mathbb{R}^3 is illustrated in Figure 1b.

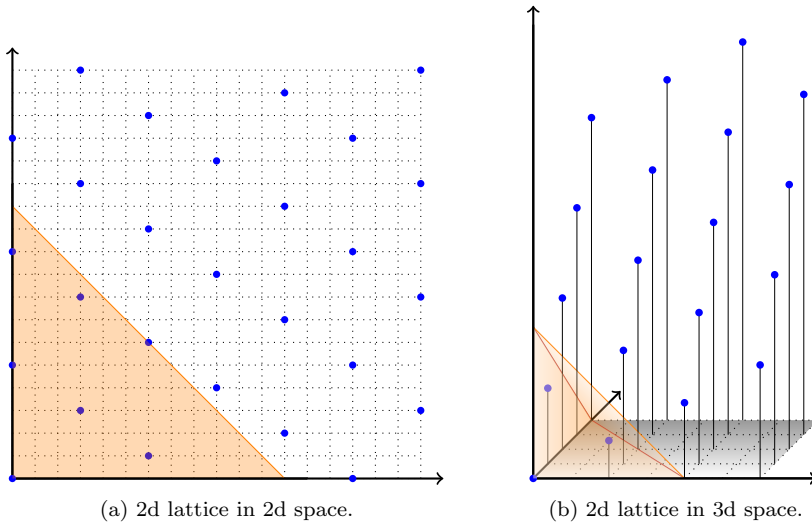


Fig. 1: Lattices in the positive orthant. The orange areas mark the dilated simplices $12S$ and $4S$ respectively.

Definition 2 (Standard Simplex) The *standard simplex* S in \mathbb{R}^m is the convex hull of the points e_0, \dots, e_m . For $d \in \mathbb{N}$, the d -dilation dS of S is the convex hull of the points de_0, \dots, de_m .

We are interested in the integer points of a dilated lattice $dS \subseteq \mathbb{R}^m$. Obviously, this set consists of all points (v_1, \dots, v_m) in \mathbb{Z}^m with $v_1, \dots, v_m \geq 0$ and $v_1 + \dots + v_m \leq d$. We can also describe it as an intersection of translated cones.

Definition 3 (Cones) A set $C \subseteq \mathbb{Z}^m$ is called a (*discrete*) *cone* if C contains the origin and we have that for all $u, v \in C$ and for all $\alpha, \beta \in \mathbb{N}$, the linear combination $\alpha u + \beta v$ is also an element of C . For vectors $c_1, \dots, c_n \in \mathbb{Z}^m$ we write $[c_1, \dots, c_n]$ for the smallest cone containing c_1, \dots, c_n , which we call generators of the cone. For $c \in C$, $[c]$ is called an *edge* of C if there exists a hyperplane $H \subset \mathbb{R}^m$ with $H \cap C \subseteq [c]$. We call edges of the form $[e_i]$ or

$[-e_i]$, $i \in \{1, \dots, m\}$ *straight*, while all other edges are called *slanted*. For $i \in \{0, \dots, m\}$, we define the i th *corner cone* C_i of the standard simplex as $[e_0 - e_i, e_1 - e_i, \dots, e_m - e_i] \subseteq \mathbb{Z}^m$.

Subsequently, we will only be concerned with finitely generated cones. We can therefore assume that a cone C is always given as a finite set of points c_i , such that for each i , $[c_i]$ is an edge of C , and for $j \neq i$ we have $[c_i] \neq [c_j]$.

The standard simplex in \mathbb{R}^m has $m + 1$ distinct corner cones C_0, \dots, C_m , and the set of all integer points in dS , $d \in \mathbb{N}$, is equal to the intersection $\bigcap_{i=0}^m (C_i - de_i)$, as illustrated for dimension 2 in Figure 2.

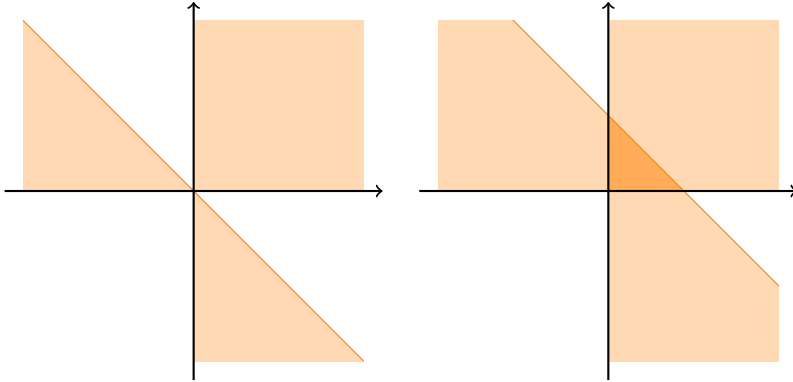


Fig. 2: Corner cones of the standard simplex and the intersection of the translated cones in \mathbb{R}^2 .

As we outlined in the earlier sections, we look for integer points in dS that are not connected to any other integer points in dS via a given lattice L . The next definition formalizes this idea not only for simplices but general subsets of \mathbb{R}^m .

Definition 4 (Lonely Points) Let $L \subseteq \mathbb{Z}^m$ be a lattice. We define the equivalence relation \sim on \mathbb{Z}^m as $u \sim v \Leftrightarrow u - v \in L$. Let A be an arbitrary subset of \mathbb{R}^m . A point $v \in A$ is called *lonely* (with respect to L), if $v \in \mathbb{Z}^m$ and there is no $\tilde{v} \in (A \cap \mathbb{Z}^m) \setminus \{v\}$ such that $\tilde{v} - v \in L$. We write $\text{lonely}_L(A)$ for the set of lonely points in A and $\#\text{lonely}_L(A) \in \mathbb{N} \cup \{\infty\}$ for the number of lonely points in A .

Example 3 We give two examples of lattices where the number of lonely points in dS does not grow indefinitely with d .

1. For $L = \langle \begin{pmatrix} 2 \\ -3 \end{pmatrix} \rangle \subseteq \mathbb{Z}^2$ there are 9 lonely points in all dS for all $d \geq 4$ (Figure 3, left).
2. For $L = \langle \begin{pmatrix} 1 \\ 1 \end{pmatrix} \rangle \subseteq \mathbb{Z}^2$ there are 4 lonely points in all dS for all $d \geq 2$. (Figure 3, right).

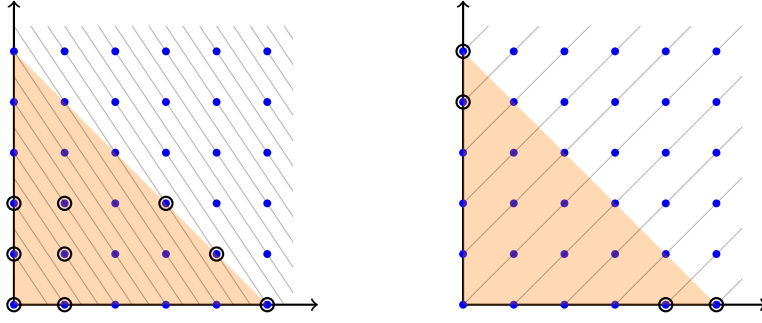


Fig. 3: Illustration of Example 3. Lonely points are encircled.

It is easy to show that there is no lattice (other than $\{0\}$) such that the number of lonely points in dS grows indefinitely with d .

Example 4 Let $L \subseteq \mathbb{Z}^4$ be the lattice generated by the vectors $(2, 0, -1, 0)$ and $(1, 1, 0, -1)$. Then there are infinitely many lonely points in any corner cone. For example, for each $i = 0, \dots, 4$, all vectors of the form $(0, n, 0, 0) - de_i$ with $d \geq n \geq 0$ are lonely in C_i .

Our goal is to count the lonely points in a dilated simplex. As we will use the translated corner cones to characterize the points inside of a dilated simplex, we want to make sure that lonely points stay lonely after any translation.

Lemma 1 *Let $L \subset \mathbb{Z}^m$ be a lattice and let $v \in A \subseteq \mathbb{R}^m$. If $v \in \text{lonely}_L(A)$, then $v + t \in \text{lonely}_L(A + t)$ for any $t \in \mathbb{Z}^m$.*

Proof Suppose $v + t \notin \text{lonely}_L(A + t)$. Then there exists a $\tilde{v} \in A$ such that $(v + t) \sim (\tilde{v} + t)$. It follows that $v - \tilde{v} = (v + t) - (\tilde{v} + t) \in L$, so $v \sim \tilde{v}$. \square

4 Counting and Identifying Lonely Points

In this section we develop algorithms for deciding whether in a given setting the number of lonely points is finite or infinite, as well as an algorithm which in the finite case determines how many lonely points there are. First we characterize loneliness of points in cones, and then we relate the loneliness of points in a dilated cone dS to the loneliness of points in its corner cones.

Lemma 2 *Let $L \subseteq \mathbb{Z}^m$ be a lattice and $C \subseteq \mathbb{Z}^m$ be a cone.*

1. *If C has any lonely points, then 0 is one of them.*
2. *C has lonely points if and only if $L \cap C = \{0\}$.*
3. *If $u \in C$ is not lonely, then also $u + v$ is not lonely for any $v \in C$.*

Proof 1. If 0 is not lonely, it is equivalent to some other point of C , say to $u \neq 0$. Then $u = u - 0 \in L$. Let v be an arbitrary element of C . Since $u \in C$, we have $v + u \in C$, and since v and $v + u$ are equivalent, v is not lonely.

2. If C has lonely points, then, by the previous item, 0 is one of them, hence $L \cap C = \{0\}$. For the other direction, if $L \cap C = \{0\}$, then 0 is lonely.
3. If u is not lonely, then there exists $\tilde{u} \in C \setminus \{u\}$ with $u \sim \tilde{u}$. Then also $u + v \sim \tilde{u} + v$, and since $\tilde{u} + v$ is in C and different from $u + v$, the claim follows. \square

Proposition 1 *Let $L \subseteq \mathbb{Z}^m$ be a lattice and $C = [c_1, \dots, c_n] \subseteq \mathbb{Z}^m$ be a cone.*

1. *If C has infinitely many lonely points, then there is an $i \in \{1, \dots, n\}$ such that all points in $[c_i]$ are lonely in C .*
2. *Let $i \in \{1, \dots, n\}$. Then all points in $[c_i]$ are lonely in C if and only if $L \cap C = \{0\}$ and $(L + \langle c_i \rangle) \cap C = [c_i]$.*

Proof 1. Suppose to the contrary all edges $[c_i]$ contain a nonlonely point, say $\alpha_1 c_1, \dots, \alpha_n c_n$ are not lonely for certain positive integers $\alpha_1, \dots, \alpha_n$. By part 3 of Lemma 2, all points $\beta_1 c_1 + \dots + \beta_n c_n$ with $\beta_1 \geq \alpha_1, \dots, \beta_n \geq \alpha_n$ are not lonely. Thus there remain only finitely many candidates for lonely points.

2. “ \Rightarrow ” If all points in $[c_i]$ are lonely, then C has lonely points, so $L \cap C = \{0\}$ by part 2 of Lemma 2. It remains to show that $(L + \langle c_i \rangle) \cap C = [c_i]$. The direction “ \supseteq ” is clear. To show “ \subseteq ”, let $v \in (L + \langle c_i \rangle) \cap C$, say $v = \ell + \alpha c_i \in C$ for some nonzero $\ell \in L$ and $\alpha \in \mathbb{Z}$. If $\alpha > 0$, then $v \sim \alpha c_i$, in contradiction to the loneliness of αc_i . Otherwise, for $\alpha \leq 0$, we have $\ell = v + (-\alpha)c_i \in C$, a contradiction to $L \cap C = \{0\}$.

“ \Leftarrow ” Assume $u = \alpha_i c_i$ is not lonely, say $u \sim v$ for some $v \in C \setminus \{u\}$. Then $u - v \in L$ implies $v \in L + \langle c_i \rangle$, so $v \in [c_i]$, say $v = \beta_i c_i$ for some $\beta \in \mathbb{N} \setminus \{\alpha_i\}$. But then $0 \neq u - v \in L \cap C = \{0\}$, a contradiction. \square

The conditions of Proposition 1 give rise to the following algorithm for deciding whether a cone contains infinitely many lonely points.

Algorithm 1 (*hasInfinitelyManyLonelyPoints*)

Input: a lattice $L \subseteq \mathbb{Z}^m$, a cone $C = [c_1, \dots, c_n] \subseteq \mathbb{Z}^m$

Output: true or false, depending on whether C contains infinitely many lonely points

- 1 if $L \cap C \neq \{0\}$ then return false
- 2 for $i = 1, \dots, n$, do:
- 3 if $(L + \langle c_i \rangle) \cap C = [c_i]$ then return true
- 4 return false

The tests in lines 1 and 3 can be performed using integer linear programming [15]. If $L = \langle \ell_1, \dots, \ell_k \rangle = [\ell_1, \dots, \ell_k, -\ell_1, \dots, -\ell_k]$, we can find nonnegative integers $\alpha_1, \dots, \alpha_k, \alpha_{-1}, \dots, \alpha_{-k}, \beta_1, \dots, \beta_n$ such that

$$(\alpha_1 - \alpha_{-1})\ell_1 + \dots + (\alpha_k - \alpha_{-k})\ell_k = \beta_1 c_1 + \dots + \beta_n c_n$$

and such that $\beta_1 + \dots + \beta_n$ is maximized. We have $L \cap C = \{0\}$ if and only if the optimal solution is $\beta_1 = \dots = \beta_n = 0$.

Similarly, in order to check whether $(L + \langle c_i \rangle) \cap C = [c_i]$, we can find nonnegative integers $\alpha_1, \dots, \alpha_k, \alpha_{-1}, \dots, \alpha_{-k}, \gamma_1, \gamma_{-1}, \beta_1, \dots, \beta_n$ such that

$$(\alpha_1 - \alpha_{-1})\ell_1 + \dots + (\alpha_k - \alpha_{-k})\ell_k + (\gamma_1 - \gamma_{-1})c_i = \beta_1 c_1 + \dots + \beta_n c_n$$

and $\beta_1 + \dots + \beta_{i-1} + \beta_{i+1} + \dots + \beta_n$ is maximized. If the intersection $[c_i] \cap [c_1, \dots, c_{i-1}, c_{i+1}, \dots, c_n]$ only contains 0, then $(L + \langle c_i \rangle) \cap C$ is contained in $[c_i]$ if and only if the optimal solution is $\beta_1 = \dots = \beta_n = 0$. In our setting, we can always assume that c_1, \dots, c_n are linearly independent over \mathbb{Q} , and in this case, the condition $[c_i] \cap [c_1, \dots, c_{i-1}, c_{i+1}, \dots, c_n] = \{0\}$ is always satisfied.

When there are only finitely many lonely points, we can next determine how many there are. Part 3 of Lemma 2 says that when some $v \in C$ is not lonely, then no point in the translated cone $v + C$ is lonely either. It follows from Dickson's lemma ([3], see also Lemma 4 of [1]) that the set of nonlonely points in C is in fact a finite union of such translated cones $v + C$, quite similar to the leading-term ideals in Gröbner basis theory [3,6,4]. Inspired by the FGLM-algorithm from that theory [8,6], we arrive at the following algorithm for counting the number of lonely points in a cone.

Algorithm 2 (*numberOfLonelyPoints*)

Input: a lattice $L \subseteq \mathbb{Z}^m$ and a cone $C = [c_1, \dots, c_n] \subseteq \mathbb{Z}^m$

Output: $\# \text{lonely}_L(C)$

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1  if  $\# \text{lonely}_L(C) = \infty$ , return  $\infty$   (( using Algorithm 1 ))
2  if 0 is not lonely, return 0
3  todo =  $\{e_1, \dots, e_n\} \subseteq \mathbb{R}^n$   (( list of unit vectors of length n ))
4  B =  $\emptyset$   (( collected nonlonely points ))
5  npoints = 1  (( number of lonely points seen so far ))
6  while |todo| > 0, do:
7     select an element  $v = (v_1, \dots, v_n)$  with  $\|v\|_1$  minimal from todo
8     todo = todo  $\setminus \{v\}$ 
9     if  $v_1 c_1 + \dots + v_n c_n$  is a lonely point, then:
10        npoints = npoints + 1
11        for  $i = 1, \dots, n$ , do:
12           if  $\forall b \in B : v_1 c_1 + \dots + v_n c_n + c_i \notin b + C$ , then
13              todo = todo  $\cup \{v + e_i\}$ 
14        else  (( v is not lonely ))
15           B = B  $\cup \{v_1 c_1 + \dots + v_n c_n\}$ 
16  return npoints

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Three aspects need to be discussed in order to justify this algorithm: (1) that all indicated operations can be performed algorithmically, (2) that it returns the correct output, and (3) that it terminates for every input. Concerning the first point, the only questionable steps are the checks in steps 2 and 9 whether a given point is lonely. In order for v to be not lonely, there must be integers $\alpha_1, \dots, \alpha_k$, not all zero, such that $v + \alpha_1 \ell_1 + \dots + \alpha_k \ell_k$ also

belongs to C , where ℓ_1, \dots, ℓ_k are generators of L . Whether such integers exist can be determined with linear programming.

For the correctness, observe first that the output $npoints$ is a lower bound on the number of lonely points, because the counter is only incremented when we have found a new lonely point v . Since we always consider the candidate of least norm and in line 13 always add elements of larger norm to the todo-list, it is excluded that we count the same point more than once. In order to see that the output is also an upper bound, observe that part 3 of Lemma 2 implies that when b is not lonely, then all the points in $b + C$ are not lonely either, so it is fair to exclude them from consideration in step 12. Since all other points will be considered, there is no danger of undercounting. This establishes the correctness.

Finally, for justifying the termination, observe that the number of iterations of the main loop is bounded by the number of lonely points plus the number of points that are not lonely but also not contained in a translated cone $b + C$ where b is a nonlonely point discovered earlier. By line 1, the number of lonely points is finite when the algorithm reaches the main loop, and we have already argued above that the number of nonlonely points not contained in a translated cone rooted at an earlier discovered nonlonely point is finite as well.

Example 5 Consider the lattice $L = \langle \begin{pmatrix} 2 \\ -3 \end{pmatrix} \rangle \subseteq \mathbb{Z}^2$ and the cone $C = [e_1, e_2] \subseteq \mathbb{Z}^2$. This cone is the corner cone C_0 in the situation considered in Example 3, part 1 and depicted in Figure 3. Algorithm 2 identifies the lonely points of C as follows.

iteration	v	todo	B	$npoints$	comment
0		$\{\begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}\}$	\emptyset	1	initialization
1	$\begin{pmatrix} 1 \\ 0 \end{pmatrix}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}\}$	\emptyset	2	v is lonely
2	$\begin{pmatrix} 0 \\ 1 \end{pmatrix}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 2 \end{pmatrix}\}$	\emptyset	3	v is lonely
3	$\begin{pmatrix} 2 \\ 0 \end{pmatrix}$	$\{\begin{pmatrix} 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 2 \end{pmatrix}\}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}\}$	3	v is not lonely
4	$\begin{pmatrix} 1 \\ 1 \end{pmatrix}$	$\{\begin{pmatrix} 1 \\ 2 \end{pmatrix}, \begin{pmatrix} 0 \\ 2 \end{pmatrix}\}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}\}$	4	v is lonely
5	$\begin{pmatrix} 0 \\ 2 \end{pmatrix}$	$\{\begin{pmatrix} 1 \\ 2 \end{pmatrix}, \begin{pmatrix} 0 \\ 3 \end{pmatrix}\}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}\}$	5	v is lonely
6	$\begin{pmatrix} 1 \\ 2 \end{pmatrix}$	$\{\begin{pmatrix} 1 \\ 3 \end{pmatrix}, \begin{pmatrix} 0 \\ 3 \end{pmatrix}\}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}\}$	6	v is lonely
7	$\begin{pmatrix} 0 \\ 3 \end{pmatrix}$	$\{\begin{pmatrix} 1 \\ 3 \end{pmatrix}\}$	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 3 \end{pmatrix}\}$	6	v is not lonely
8	$\begin{pmatrix} 1 \\ 3 \end{pmatrix}$	\emptyset	$\{\begin{pmatrix} 2 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 3 \end{pmatrix}, \begin{pmatrix} 1 \\ 3 \end{pmatrix}\}$	6	v is not lonely

The next proposition connects the lonely points in a simplex to the lonely points in its corner cones.

Proposition 2 *Let $L \subseteq \mathbb{Z}^m$ be a lattice and let $S \subseteq \mathbb{R}^m$ be the standard simplex.*

1. *A corner de_i of dS is lonely for all sufficiently large $d \in \mathbb{N}$ iff 0 is a lonely point of the corresponding corner cone C_i .*
2. *$\forall d \in \mathbb{N} : \text{lonely}_L(dS) \supseteq \bigcup \{v - de_i \mid \exists i > 0 : v \in \text{lonely}_L(C_i)\} \cap \mathbb{N}^m$.*
3. *The following are equivalent:*

- (a) $\forall i : \# \text{lonely}_L(C_i) = \infty.$
- (b) $\exists i : \# \text{lonely}_L(C_i) = \infty,$
- (c) $\forall r \in \mathbb{N} \exists d \in \mathbb{N} : \# \text{lonely}_L(dS) > r,$

Proof 1. Let de_i be a corner of dS , and suppose d is large.

“ \Rightarrow ” We show: if 0 is not a lonely point of the corner cone $C_i = [e_0 - e_i, \dots, e_m - e_i]$, then de_i is not a lonely point of dS . If 0 is not a lonely point of the corner cone, the corner cone contains some nonzero element of L , say $\ell = \alpha_0(e_0 - e_i) + \dots + \alpha_m(e_m - e_i) \in L$ for certain $\alpha_0, \dots, \alpha_m \in \mathbb{N}$. Assuming, as we may, that $d > \alpha_0 + \dots + \alpha_m$, we have that $de_i + \ell$ is an interior point of dS which is equivalent to de_i , proving that de_i is not lonely.

“ \Leftarrow ” We show: if de_i is not a lonely point of dS , then 0 is not a lonely point of the corner cone. Indeed, suppose that de_i is equivalent to another point v of dS , say to $v = \beta_1 e_1 + \dots + \beta_m e_m$ for some $\beta_1, \dots, \beta_m \geq 0$ whose sum is at most d . Then $v - de_i = \beta_1(e_1 - e_i) + \dots + \beta_m(e_m - e_i) + (d - \sum_j \beta_j)(e_0 - e_i)$ belongs to the i th corner cone, so 0 is not a lonely point of that cone.

2. Denote the set on the right hand side by A_d . Then $A_d \subset \text{lonely}_L(dS)$ holds for any d : If $v - de_i \in \mathbb{N}^m$ is such that v is lonely in C_i , then by Lemma 1 $v - de_i$ is lonely in $C_i - de_i$, which contains dS .
3. “(a) \Rightarrow (b)” is trivial.

“(b) \Rightarrow (c)” is an immediate consequence of part 2.

“(c) \Rightarrow (a)” Suppose that $\text{lonely}_L(C_i)$ only contains finitely many elements for some corner cone $C_i = [c_1, \dots, c_n]$. Then, by part 1 of Proposition 1 there exists a d' such that for every edge $[c_j]$ in C_i the point $d'c_j$ is not lonely. For each such edge we let d_j be the minimal euclidean distance of $d'c_j$ to some other element in C_i equivalent to $d'c_j$. Then any point $v = \sum \alpha_j c_j$ in C_i is equivalent to some point in distance d_j if $\alpha_j \geq d'$ for some j . Setting d to be the maximum of the d_j this means that every such v is equivalent to some point in distance $\leq d$. Then a point v in dS for $\tilde{d} \geq d$ is lonely only if the coordinates of $v - \tilde{d}e_i$ with respect to the generators c_j of the i th corner cone are bounded by d , leaving only finitely many possible values for $v - \tilde{d}e_i$. \square

For a specific $d \in \mathbb{N}$, there are only finitely many points in dS , and for each of them, we can decide whether it is lonely in a similar way as described above for a given point in a cone. The issue reduces to a linear programming question. What we are interested in is how far the number of lonely points can grow as d increases. Proposition 2 says that the lonely points in dS for sufficiently large d are essentially the lonely points of the corner cones. When a cone has only finitely many lonely points, they are all clustered near the apex, so as soon as d is sufficiently large, the number of lonely points in the dilated simplex dS is exactly the sum of the number of lonely points in its corner cones. When at least one corner cone has infinitely many lonely points, then the number of lonely points in dS is unbounded as d goes to infinity. In summary, we obtain the following algorithm.

Algorithm 3 (*ultimateNumberOfLonelyPoints*)**Input:** a lattice $L \subseteq \mathbb{Z}^m$ **Output:** $\lim_{d \rightarrow \infty} \# \text{lonely}_L(dS)$

- 1 $s = 0$
- 2 for all $i \in \{0, \dots, m\}$, do:
- 3 $C = [e_0 - e_i, \dots, e_m - e_i]$ ((consider the i th corner cone))
- 4 $s = s + \# \text{lonely}_L(C)$ ((use Algorithm 2))
- 5 return s

We have implemented the algorithms described in this section in Mathematica. The code is available on the personal website of the second author.

5 Lonely Points for Small Lattices

It is clear that all integer points in dS are lonely when $L = \{0\}$ and that there are no lonely points when $L = \mathbb{Z}^m$. More generally, geometric intuition suggests that there should be more lonely points when L is “small”. The main result of the present section makes this intuition quantitative. We show that whenever the dimension of L is less than a certain constant multiple of the ambient dimension m , then there is a corner cone which satisfies the conditions of part 2 of Proposition 1 and thus has infinitely many lonely points.

In the subsequent proofs we make use of sign vectors and equations. The possible components of a sign vector are $+$, $-$, \oplus , \ominus or 0 . We can assign a sign vector s to a given $v \in \mathbb{R}^m$ in the following way. If the i th component of v is nonnegative, then the i th component of s is $+$ or \oplus . If the i th component of v is nonpositive, then the i th component of s is $-$ or \ominus . If a component of v is zero, then the corresponding component of s can be 0 , $+$, $-$, \oplus or \ominus . A component of s is \oplus or \ominus only if the absolute value of the corresponding component of v is greater than or equal to the sum of the absolute values of all other components. With these rules, any equation $s_1 + \dots + s_k = s$ of sign vectors s_1, \dots, s_k, s is a valid equation if there are vectors $v_1, \dots, v_k, v \in \mathbb{R}^m$, such that $v_1 + \dots + v_k = v$ and for each $i = 1, \dots, k$, s_i is a valid sign vector for v_i and s is a valid sign vector for v .

Example 6 For the equation

$$\begin{pmatrix} -2 \\ 1 \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ -1 \\ 0 \end{pmatrix} + \begin{pmatrix} 1 \\ 1 \\ -1 \end{pmatrix} = \begin{pmatrix} -1 \\ 1 \\ -1 \end{pmatrix},$$

two valid sign equations are

$$\begin{pmatrix} \ominus \\ + \\ - \end{pmatrix} + \begin{pmatrix} + \\ \ominus \\ - \end{pmatrix} + \begin{pmatrix} + \\ + \\ - \end{pmatrix} = \begin{pmatrix} - \\ + \\ - \end{pmatrix}, \text{ and } \begin{pmatrix} - \\ + \\ 0 \end{pmatrix} + \begin{pmatrix} 0 \\ - \\ 0 \end{pmatrix} + \begin{pmatrix} + \\ + \\ - \end{pmatrix} = \begin{pmatrix} - \\ + \\ - \end{pmatrix}.$$

We use a shorthand matrix notation

$$[s_1 \ s_2 \ \cdots \ s_k] = s,$$

for the sign equation $s_1 + \cdots + s_k = s$, with the square brackets indicating that the columns of the matrix are summed up to obtain the right hand side. To further shorten notation, we use $\boxed{\oplus}$ and $\boxed{\ominus}$ for nonempty square blocks of the form

$$\begin{array}{ccc} \oplus - \cdots \cdots \cdots - - & & \ominus + \cdots \cdots \cdots + + \\ - \oplus & & + \ominus & & + + \\ - - \ddots & & - - & & + + \ddots & & + + \\ \vdots \vdots \ddots & & \vdots \vdots & \text{and} & \vdots \vdots \ddots & & \vdots \vdots \\ - - & & \ddots - - & & + + & & \ddots + + \\ - - & & \oplus - & & + + & & \ominus + \\ - - \cdots \cdots \cdots - \oplus & & & & + + \cdots \cdots \cdots + \ominus \end{array}$$

respectively, where the number of rows/columns is either clear from the context or irrelevant. Similarly we use $\boxed{+}$, $\boxed{-}$, and $\boxed{0}$ for blocks that only contain $+$, $-$, or 0 respectively, with the difference that these blocks do neither have to be square blocks nor nonempty.

Example 7 The first sign equation in Example 6 can be written as

$$\boxed{\ominus} + \begin{pmatrix} + \\ + \\ - \end{pmatrix} = \begin{pmatrix} - \\ + \\ - \end{pmatrix}.$$

For any vector v in \mathbb{R}^m we define the *balance* $\tau(v)$ of v to be the sum of the components of v . The balance of a vector v with only nonnegative components is equal to $\|v\|_1$. For any slanted edge $[c]$ of a corner cone, c is the difference of two unit vectors, and thus $\tau(c) = 0$. For straight edges we have $\tau(c) = \tau(\pm e_i) = \pm 1$.

Definition 5 (Visible Vectors) We call a vector $v \in \mathbb{R}^m$ *i-visible*, if

$$(+, \dots, +, \ominus, +, \dots, +)$$

is a valid sign vector for v , where \ominus is at the i th position.

The definition is motivated by corner cones. For $i > 0$, a vector is *i-visible* iff it belongs to C_i . An *i-visible* vector v has nonpositive balance $\tau(v) \leq 0$.

Lemma 3 Let $k \in \{1, \dots, m\}$ and let $v_1, \dots, v_k, v \in \mathbb{R}^m$ be such that $v_1 + v_2 + \dots + v_k = v$ and that each v_i lies in some corner cone. Suppose that there is an associated sign equation and indices r_1, \dots, r_k such that a valid sign equation projected to rows r_1, \dots, r_k is of the form

$$\boxed{\ominus} = \left(\begin{array}{c} \boxed{-} \\ \boxed{+} \end{array} \right).$$

Then for every $j \in \{r_1, \dots, r_k\}$, the j th component of v is zero, and for every $j \in \{1, \dots, m\} \setminus \{r_1, \dots, r_k\}$, the j th component of v_i is zero for every i .

Proof Let $\pi: \mathbb{R}^m \rightarrow \mathbb{R}^n$ be the projection on the components with indices r_1, \dots, r_k , and $\bar{\pi}$ the projection on the complementary components. The sign equation implies that $\tau(\pi(v_i)) \leq 0$ for all v_i . It follows that $\tau(\pi(v))$ has to be less than or equal to 0 as well. As $\pi(v)$ only contains nonnegative entries, this is only possible if $\pi(v)$ is the zero vector. This shows the first part and also implies the equation

$$\tau(\pi(v_1)) + \tau(\pi(v_2)) + \dots + \tau(\pi(v_k)) = 0.$$

Since no summand on the left hand side is strictly positive, all the $\tau(\pi(v_i))$ have to be equal to 0. As every v_i lies in some corner cone, and their negative components only have indices contained in $\{r_1, \dots, r_k\}$, we get that all the $\bar{\pi}(v_i)$ only have nonpositive components. Now it follows that all $\bar{\pi}(v_i)$ are equal to zero, since

$$0 \geq \tau(\bar{\pi}(v_i)) = \tau(\bar{\pi}(v_i)) + \tau(\pi(v_i)) = \tau(v_i) \geq 0. \quad \square$$

Remark 1 Clearly, if a v_1, \dots, v_k, v with $v_1 + \dots + v_k = v$ are such that a valid sign equation contains rows of the form

$$\begin{bmatrix} \pm \\ \pm \\ \pm \end{bmatrix} = \begin{bmatrix} \mp \\ \mp \\ \mp \end{bmatrix},$$

then v and the v_i can only contain zero entries at the corresponding indices.

Proposition 3 *Let $v_1, \dots, v_m \in \mathbb{R}^m \setminus \{0\}$ be such that each v_i is i -visible. If no subspace of $V := \langle v_1, \dots, v_m \rangle$ can be decomposed into a direct sum of more than one nonzero vector spaces, then $\dim(V) = m - 1$.*

Proof If V is of dimension m , it can be decomposed into the direct sum of m nonzero vector spaces. Suppose that $\dim(V) < m - 1$, and, without loss of generality, that v_1, \dots, v_{m-2} generate V , i.e. there are $\alpha_1, \dots, \alpha_{m-2}$ and $\beta_1, \dots, \beta_{m-2} \in \mathbb{R}$ with $v_{m-1} = \sum_{i < m-1} \alpha_i v_i$ and $v_m = \sum_{i < m-1} \beta_i v_i$. For these we get corresponding sign equations

$$\pm \begin{pmatrix} \ominus \\ + \\ \vdots \\ + \\ + \\ + \\ + \end{pmatrix} \pm \dots \pm \begin{pmatrix} + \\ + \\ \vdots \\ + \\ \ominus \\ + \\ + \end{pmatrix} = \begin{pmatrix} + \\ + \\ \vdots \\ + \\ + \\ \ominus \\ + \end{pmatrix}, \quad (2)$$

$$\pm \begin{pmatrix} \ominus \\ + \\ \vdots \\ + \\ + \\ + \\ + \end{pmatrix} \pm \dots \pm \begin{pmatrix} + \\ + \\ \vdots \\ + \\ \ominus \\ + \\ + \end{pmatrix} = \begin{pmatrix} + \\ + \\ \vdots \\ + \\ + \\ + \\ \ominus \end{pmatrix}, \quad (3)$$

where the \pm reflect the fact that the α_i and β_i can be positive or negative. We show that there is no combination of signs for the α_i and β_i such that both equalities hold, unless V can be decomposed into a direct sum. We first look at (3). From the last row we see that at least one β_i has to be strictly negative, as no v_i on the left hand side has a negative entry at index m , and v_m is nonzero. So we can split the vectors into two groups: those with positive β_i and those with strictly negative β_i . After changing the summation order and reorganizing the rows if necessary, Equation (3) becomes

$$\begin{bmatrix} \ominus \\ \boxed{+} \\ + \cdots + \\ + \cdots + \end{bmatrix} + \begin{bmatrix} \boxed{-} \\ \oplus \\ - \cdots - \\ - \cdots - \end{bmatrix} = \begin{pmatrix} \boxed{+} \\ \boxed{+} \\ + \\ \ominus \end{pmatrix}.$$

Suppose at least one β_i is strictly positive. Then Lemma 3 implies that some components have to be zero, and we get a block diagonal form

$$\begin{bmatrix} \ominus \\ \boxed{0} \\ 0 \cdots 0 \\ 0 \cdots 0 \end{bmatrix} + \begin{bmatrix} \boxed{0} \\ \oplus \\ 0 \cdots 0 \\ - \cdots - \end{bmatrix} = \begin{pmatrix} \boxed{+} \\ \boxed{+} \\ 0 \\ \ominus \end{pmatrix}.$$

Thus, the v_i appearing with a nonzero coefficient in (2) span a vector space that can be decomposed into a direct sum if at least one β_i is strictly positive. Otherwise, with the analogous reasoning for v_{m-1} , we can suppose that all the α_i and β_i are nonpositive, and conclude that the sign equations for v_{m-1} and v_m are of the form

$$\begin{bmatrix} \oplus \\ - \cdots - \\ - \cdots - \end{bmatrix} = \begin{pmatrix} \boxed{+} \\ \ominus \\ + \end{pmatrix}, \quad (4)$$

$$\begin{bmatrix} \oplus \\ - \cdots - \\ - \cdots - \end{bmatrix} = \begin{pmatrix} \boxed{+} \\ + \\ \ominus \end{pmatrix}. \quad (5)$$

If all α_i (implicitly used in (4)) were nonzero, then the last row in (4) implies that the last components of all the v_i would have to be zero, which is incompatible with the last row in (5). Again with the analogous reasoning for the β_i we see that not all α_i and not all β_i are nonzero. As before we split the vectors on the left hand side of each equation into two blocks: vectors that appear with a nonzero coefficient in only one of the equations and vectors that are shared in both equations with nonzero coefficients, which gives, after reordering the

rows and summands if necessary:

$$\begin{aligned} \begin{bmatrix} \boxed{\oplus} \\ \boxed{-} \\ \boxed{-} \\ \dots \\ \boxed{-} \end{bmatrix} + \begin{bmatrix} \boxed{-} \\ \boxed{\oplus} \\ \boxed{-} \\ \dots \\ \boxed{-} \end{bmatrix} &= \begin{pmatrix} \boxed{-} \\ \boxed{+} \\ \boxed{+} \\ \boxed{+} \\ \ominus \\ \boxed{+} \end{pmatrix}, \\ \underbrace{\begin{bmatrix} \boxed{-} \\ \boxed{\oplus} \\ \boxed{-} \\ \dots \\ \boxed{-} \end{bmatrix}}_{\text{shared}} + \begin{bmatrix} \boxed{-} \\ \boxed{\oplus} \\ \boxed{-} \\ \dots \\ \boxed{-} \end{bmatrix} &= \begin{pmatrix} \boxed{-} \\ \boxed{+} \\ \boxed{+} \\ \boxed{+} \\ \boxed{+} \\ \ominus \end{pmatrix}. \end{aligned}$$

We use Remark 1 to determine zero components in the first equation:

$$\begin{bmatrix} \boxed{\oplus} \\ \boxed{-} \\ \boxed{0} \\ \dots \\ \boxed{0} \end{bmatrix} + \underbrace{\begin{bmatrix} \boxed{-} \\ \boxed{\oplus} \\ \boxed{0} \\ \dots \\ \boxed{0} \end{bmatrix}}_{\text{shared}} = \begin{pmatrix} \boxed{-} \\ \boxed{+} \\ \boxed{+} \\ \boxed{0} \\ \ominus \\ \boxed{0} \end{pmatrix}.$$

Then, doing the same for the second equation, and using the fact that we already know some zero components in the shared vectors, we get:

$$\underbrace{\begin{bmatrix} \boxed{0} \\ \boxed{\oplus} \\ \boxed{0} \\ \dots \\ \boxed{0} \end{bmatrix}}_{\text{shared}} + \begin{bmatrix} \boxed{0} \\ \boxed{-} \\ \boxed{\oplus} \\ \dots \\ \boxed{-} \end{bmatrix} = \begin{pmatrix} \boxed{0} \\ \boxed{+} \\ \boxed{+} \\ \boxed{+} \\ \boxed{+} \\ \ominus \end{pmatrix}.$$

Denote the number of shared vectors by s . If s is greater than 0, we look at the rows in the equation for v_{m-1} where the shared vectors are nonzero:

$$\underbrace{\begin{bmatrix} \boxed{-} \\ \boxed{-} \end{bmatrix}}_{s \text{ many shared}} + \underbrace{\begin{bmatrix} \boxed{\oplus} \end{bmatrix}}_{\text{shared}} = \begin{pmatrix} \boxed{-} \\ \boxed{+} \end{pmatrix}.$$

As all nonshared vectors on the left hand side only have negative components, we can bring them to the right hand side and get:

$$\underbrace{\left[\begin{array}{c} \oplus \\ \oplus \\ \oplus \\ \oplus \end{array} \right]}_{s \text{ many shared}} = \left(\begin{array}{c} [-] \\ [-] \\ [-] \\ [-] \end{array} \right).$$

Note that here, all the hidden entries of the shared vectors are zero. We can suppose that the shared vectors are linearly independent, otherwise we could replace some coefficients with zero. As they are linearly independent, however, they span the whole space \mathbb{R}^s , thus the shared vectors can be replaced by unit vectors, which leads to a decomposition of V into a direct sum of vector spaces. It remains to handle the case where there are no shared vectors in (4) and (5). In that case, certain components in (4) and (5) have to be zero:

$$\begin{aligned} \left[\begin{array}{c} \oplus \\ 0 \\ - \dots - \\ 0 \dots 0 \end{array} \right] &= \left(\begin{array}{c} [-] \\ [-] \\ 0 \\ 0 \end{array} \right), \\ \left[\begin{array}{c} 0 \\ \oplus \\ 0 \dots 0 \\ - \dots - \end{array} \right] &= \left(\begin{array}{c} [-] \\ [-] \\ 0 \\ 0 \end{array} \right). \end{aligned}$$

The zero entries on the left hand side imply that the space spanned by V can be decomposed into a direct sum. This completes the proof. \square

Corollary 1 *Let $L \subseteq \mathbb{Z}^m$, $m \geq 3$, be a lattice of dimension less than $m - 1$ such that no subspace of the vector space spanned by L can be decomposed into a direct sum of two nonzero spaces. Then there exists a corner cone C of the standard simplex such that $L \cap C = \{0\}$.*

Proof Let $v_i \in L \cap C_i$ for all $i = 1, \dots, m$ and assume they are all nonzero. Then, Proposition 3 yields:

$$\dim(L) \geq \dim\langle v_1, \dots, v_m \rangle = m - 1,$$

a contradiction. \square

Corollary 2 *Let $L = \langle v_1, \dots, v_k \rangle$ be a lattice in \mathbb{Z}^m , $m \geq 3$. If $k = \dim(L) < \frac{2}{3}m$, then there exists a corner cone C of the standard simplex such that $L \cap C = \{0\}$.*

Proof Using Corollary 1 and projecting to the relevant coordinates shows that any subset S of v_1, \dots, v_k of some cardinality s such that S cannot be decomposed into a direct sum can only contain nonzero vectors of $s + 1$ corner cones.

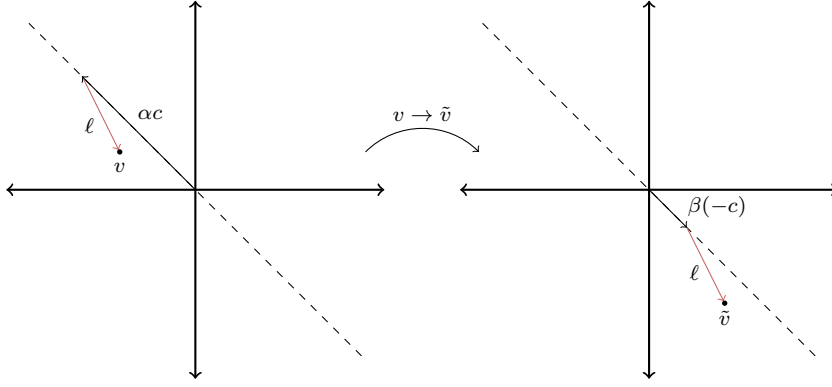


Fig. 4: Illustration of Lemma 4 in dimension 2 with $c = (-1, 1)$, $\ell = (1, -2)$, $v = (-2, 1)$, $\alpha = 3$, $\gamma = 4$, $\beta = 1$ and $\tilde{v} = (2, -3)$.

Additionally for $s = 1$, S can only contain nonzero vectors of one corner cone. In fact, if there is a $v \in \{v_1, \dots, v_k\}$ and an $i \in \{1, \dots, m\}$ such that v is i -visible, then it is immediate from the sign vector $(+, \dots, +, \ominus, +, \dots, +)$ of v that $\langle v \rangle \cap C_j = \{0\}$ for all $j \neq i$. It follows that V can be decomposed into the sum of at most $\frac{k}{2}$ many two-dimensional vector spaces, each containing nonzero vectors of 3 corner cones. \square

In order to derive a dimension bound such that both conditions in part 2 of Proposition 1 are met, we need the following lemma that allows us to construct a nonlonely point in a corner cone from a nonlonely point in a different corner cone. A geometric interpretation of the statement is given in Figure 4.

Lemma 4 *Let C_i be a corner cone, $[c]$ be a slanted edge in C_i , and let $j \in \mathbb{N}$ be such that the j th component of c is 1. If $\ell \in L$ and $\alpha \in \mathbb{N}$ are such that $v := \ell + \alpha c \in C_i$, then there exists a $\beta \in \mathbb{N}^*$ with $\ell + \beta(-c) \in C_j \setminus \{0\}$, where $[-c]$ is a slanted edge in the corner cone C_j .*

Proof By definition, the components of c are all zero except for the i th component, which is -1 , and the j th component for some $j \neq i$, which is 1. Thus $[-c]$ is a slanted edge in C_j . Set $\gamma := \max(-v_i, \alpha) + 1$, where v_i is the i th component of v . Then $\tilde{v} := v - \gamma c$ is j visible, as $\tilde{v}_i = v_i + \gamma > 0$, $\tilde{v}_k = v_k \geq 0$ for all $k \neq i, j$ and $\tilde{v}_j = v_j - \gamma \leq 0$ with

$$-\tilde{v}_j = -v_j + \gamma = -v_j - v_i + (v_i + \gamma) > -v_j + \sum_{k \neq i} v_k + \tilde{v}_i = \sum_{k \neq j} \tilde{v}_k.$$

Then, with $\beta := \gamma - \alpha \in \mathbb{N}^*$, we get $\ell + \beta(-c) = \ell + (\alpha - \gamma)c = \tilde{v} \in C_j$. \square

Theorem 1 *Let L be a lattice in \mathbb{Z}^m . If $\dim(L) < \frac{1}{3}(m-4)$, then there exists a slanted edge $[c]$ in a corner cone C such that $L \cap C = \{0\}$ and $(L + \langle c \rangle) \cap C = [c]$.*

Proof If $m \leq 4$, there is nothing to show. Suppose $m > 4$ and, without loss of generality, even. If there exist such c and C , then there is no nonzero $\ell \in L$ and no nonzero $\alpha \in \mathbb{Z}$ such that $\ell + \alpha c \in C$. Thus we can prove the claim by showing that if there exist $m(m-1)$ many nonzero $\ell_{1,2}, \ell_{1,3}, \dots, \ell_{2,1}, \ell_{2,3}, \dots, \ell_{m,m-1} \in L$ and nonzero $\alpha_{1,2}, \dots, \alpha_{m,m-1} \in \mathbb{Z}$ such that for each corner cone C_i and each slanted edge $[c_{i,j}]$ the vector $v_{i,j} := \ell_{i,j} + \alpha_{i,j}c_{i,j}$ is i -visible, then $\dim(L) \geq \frac{1}{3}(m-4)$. So suppose such $\ell_{i,j}$ and $\alpha_{i,j}$ exist. Then $(\begin{smallmatrix} + \\ + \\ - \end{smallmatrix}, \ominus, \begin{smallmatrix} + \\ - \\ - \end{smallmatrix})$ is a valid sign vector for $v_{i,j}$, with \ominus at the i th position. We first show that each $\ell_{i,j}$ is either i -visible, j -visible, or has exactly two strictly negative entries, at indices i and j . For the moment, we focus on $i=1, j=2$, allowing us to drop both indices. The reasoning for all other pairs i, j is analogous. We get the equation: $\ell + \alpha c = v$.

If $\alpha \leq 0$, we can add $-\alpha c$ to both sides of the equation, not perturbing the 1-visibility of the right hand side, which shows that ℓ is 1-visible. Otherwise, we get a sign equation with unknown entries for ℓ ,

$$\begin{pmatrix} ? \\ ? \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix} + \begin{pmatrix} \ominus \\ \oplus \\ \begin{smallmatrix} 0 \\ - \end{smallmatrix} \end{pmatrix} = \begin{pmatrix} \ominus \\ + \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix}.$$

ℓ αc v

The signs for all but two components of ℓ are immediate:

$$\begin{pmatrix} ? \\ ? \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix} + \begin{pmatrix} \ominus \\ \oplus \\ \begin{smallmatrix} 0 \\ - \end{smallmatrix} \end{pmatrix} = \begin{pmatrix} \ominus \\ + \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix}.$$

As $\tau(c) = 0$ and $\tau(v) \leq 0$, we get that $\tau(\ell) \leq 0$. Thus, if the second component of ℓ is positive, then it follows that ℓ is 1-visible with a strictly negative first component. If the second component of ℓ is negative, we can apply Lemma 4 to see that there exists a $\beta \in \mathbb{N}^*$ such that $\ell + \beta(-c)$ is 2-visible, yielding

$$\begin{pmatrix} ? \\ - \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix} + \begin{pmatrix} \oplus \\ \ominus \\ \begin{smallmatrix} 0 \\ - \end{smallmatrix} \end{pmatrix} = \begin{pmatrix} + \\ \ominus \\ \begin{smallmatrix} - \\ - \end{smallmatrix} \end{pmatrix}$$

With the same reasoning as above we can determine that ℓ is either 2-visible or its first component is strictly negative. This shows our claim for the $\ell_{i,j}$. It follows that for each pair (i, j) , the vector $\ell_{i,j}$ is such that it has a strictly negative entry at i , or j , or both. Thus we can find at least $\frac{m}{2}$ pairwise different $\ell_1, \dots, \ell_{\frac{m}{2}} \in L$ such that no two ℓ_i have a negative entry at the same index, and for each index in $\{1, \dots, m\}$, there is exactly one ℓ_i with a negative entry at that position. We now map these lattice elements to i -visible vectors, $i = 1, \dots, \frac{m}{2} - 2$, in $\mathbb{Z}^{m/2}$.

For any permutation π of $1, \dots, m$ consider the surjective linear map

$$\begin{aligned} \psi_\pi : \mathbb{R}^m &\rightarrow \mathbb{R}^{m/2}, \\ (u_1, \dots, u_m) &\mapsto (u_{\pi(1)} + u_{\pi(2)}, u_{\pi(3)} + u_{\pi(4)}, \dots, u_{\pi(m-1)} + u_{\pi(m)}). \end{aligned}$$

There are $n(m) := m!/2^{m/2}$ many such maps. We say a vector u and a map ψ_π are compatible, if:

- u is i -visible for some i , and $\psi_\pi(u) \neq 0$. If $a \in \mathbb{N}$ is such that $\pi(a) = i$, then ψ_π is $\lfloor \frac{a+1}{2} \rfloor$ -visible.
- u contains exactly two strictly negative entries at indices i and j , and there is an odd integer a such that $\pi(a) = i$ and $\pi(a+1) = j$, i.e. when applying ψ_π on v , the two negative entries are added together to give an $\frac{a+1}{2}$ -visible vector.

We now show that there exists a permutation π such that at least $\frac{m}{2} - 2$ many ℓ_i are compatible to ψ_π . In fact we can choose π such that all ℓ_i with exactly two negative entries are compatible with ψ_π , as they do not have negative entries at the same indices. This leaves us with some even number $k \geq 0$ of indices not yet considered for π and k many ℓ_i that could potentially be incompatible to such a permutation. Furthermore, there are $n(k)$ many permutations left to choose from. Each of the remaining ℓ_i is contained in a different corner cone, say C_i , and so ℓ_i is incompatible if $\psi_\pi(\ell_i) = 0$. For $k > 2$, each ℓ_i can be in the kernel of at most $n(k-2)$ many of the remaining permutations (this is the case if ℓ_i is contained in a slanted edge of a corner cone). As there are k ($k > 2$, even) many such ℓ_i , there has to be a ψ_π for which the number of i -visible ℓ_i that are mapped to zero is at most

$$\left\lfloor k \frac{n(k-2)}{n(k)} \right\rfloor = \left\lfloor \frac{2}{k-1} \right\rfloor = 0.$$

For $k = 2$, there is only one choice for π , and we could be in the situation where both of the ℓ_i have to be mapped to zero. For any such π , the images of the ℓ_i therefore contain at least $m/2 - 2$ many nonzero vectors with $m/2 - 2$ different sign patterns (after potentially reordering the rows)

$$\begin{pmatrix} \ominus \\ + \\ \vdots \\ + \\ + \\ + \end{pmatrix}, \begin{pmatrix} + \\ \ominus \\ \vdots \\ + \\ + \\ + \end{pmatrix}, \dots, \begin{pmatrix} + \\ + \\ \vdots \\ \ominus \\ + \\ + \end{pmatrix}.$$

By projecting to the first $\frac{m}{2} - 2$ coordinates and using Corollary 2, it follows that

$$\dim\langle \ell_1, \dots, \ell_{\frac{m}{2}} \rangle = \dim\langle \psi_\pi(\ell_1), \dots, \psi_\pi(\ell_{\frac{m}{2}}) \rangle \geq \frac{2}{3} \left(\frac{m}{2} - 2 \right) = \frac{1}{3}(m-4).$$

This proves the claim. \square

Without further restrictions on L , there is no analogous result for straight edges, i.e. there is no upper bound for the dimension proportional to m such that lower dimensional lattices necessarily lead to infinitely many lonely points on at least one straight edge. For any m , the lattice generated by $(1, 0, \dots, 0)$ yields only finitely many lonely points on any straight edge.

6 Conclusion and Open Questions

We translated the problem of reducing the order of a C-finite sequence to questions about which points in a dilated simplex are not connected to any other point in the simplex via a specific lattice. Our answers to these questions are in the form of algorithms that determine when the number of these points grows indefinitely with the dilation, and also compute the exact number if there are only finitely many lonely points. Furthermore we showed that if the dimension of the lattice is small enough, then the number of lonely points always grows indefinitely.

Theorem 1 is helpful for our original application to C-finite sequences, because the lattices appearing in this context are typically small. We do not know however whether the bound of Theorem 1 is tight enough to cover all cases of interest. If it is not, we can still use the Algorithms from Section 4 to see whether there are enough lonely points to derive a finite degree bound for the ansatz. At the moment, we do not know whether this is always the case.

As for extensions of our theoretical results, there are immediate questions that are rooted in discrete geometry: Can we find a closed form expression depending on d for the number of lonely points in dS for a given lattice? How many lonely points are there in more involved convex polytopes? How do linear transformations on the lattice affect lonely points? While these kinds of questions are more removed from the initial number theoretic application, their pursuit may lead to valuable insight.

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